How Slow is Quadruple Precision?

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ICERM Workshop on Variable Precision in Mathematical and Scientific Computing **Workshop Abstract:** [...] Exascale computing has also exposed the need for even greater precision than IEEE 64-bit double in some cases, because greatly magnified numerical sensitivities often mean that one can no longer be certain that results are numerically reliable. One remedy is to use IEEE 128-bit quad precision in selected portions of the computation, which is now available via software in some compilers, notably the gfortran compiler. As a single example, researchers at Stanford have had remarkable success in using quad precision in multiscale linear programming applications in biology. [...]

Plan of the Talk

- the IEEE-754 binary128 format
- a toy example: the double pendulum
- can we do better?
- conclusion and perspectives

The IEEE-754 Binary128 Format

Encoding:

sign <i>s</i>	exponent <i>e</i>	significand <i>m</i>
1 bit	15 bits	112 bits

Decoded value (except special numbers):

$$x = (-1)^s \cdot 2^{e-16383} \cdot (1 + \frac{m}{2^{112}})$$

Smallest absolute value: $x_{\rm min} \approx 6.5 \cdot 10^{-4966}$ Largest absolute value: $x_{\rm max} \approx 5.9 \cdot 10^{4931}$ Accuracy about 34 decimal digits. Currently, only the IBM Power9 supports binary128 in hardware.

Several compilers/libraries support binary128 in software:

- GNU libc/libquadmath (_Float128);
- Intel Math library (_Quad);
- Berkeley's SoftFloat by John Hauser;
- Oracle Studio;
- ASquadmath by Alexei Sibidanov (not publicly available).

Example: the Double Pendulum



 θ_1 (θ_2): angle of the 1st (2nd) pendulum wrt the vertical axis

$$u = \theta_2'^2 \ell_2 + \theta_1'^2 \ell_1 \cos(\theta_1 - \theta_2)$$

$$v = g(2m_1 + m_2) \sin \theta_1$$

$$w = m_2 g \sin(\theta_1 - 2\theta_2)$$

$$x = \theta_1'^2 \ell_1 (m_1 + m_2)$$

$$y = g(m_1 + m_2) \cos \theta_1$$

$$z = \theta_2'^2 \ell_2 m_2 \cos(\theta_1 - \theta_2)$$

$$d = 2m_1 + m_2 - m_2 \cos(2\theta_1 - 2\theta_2)$$

$$\theta_1'' = \frac{-v - w - 2 \sin(\theta_1 - \theta_2) m_2 u}{\ell_1 d}$$

$$\theta_2'' = \frac{2 \sin(\theta_1 - \theta_2) (x + y + z)}{\ell_2 d}$$

How Slow is Quadruple Precision?

Testing Framework

Pendulum lengths: $\ell_1 = \ell_2 = 1$. Masses: $m_1 = m_2 = 2$. Acceleration due to gravity: g = 9.81. Initial conditions: $\theta_1(0) = \theta_2(0) = \pi/2$, with $\theta'_1(0) = \theta'_2(0) = 0$. Integration time: 20 seconds. Method: Euler's scheme, with 50 000 steps per second (h = 1/50000):

$$egin{array}{rcl} heta_i'(t+h) &=& heta_i'(t)+h heta_i''(t)\ heta_i(t+h) &=& heta_i(t)+h heta_i'(t) \end{array}$$

Source code: http://www.loria.fr/~zimmerma/double_pendulum.html

How Slow is Quadruple Precision?

miriel038.plafrim.cluster: Intel Xeon E5-2680, 2.50GHz Ratio to the reference time of glibc/double (220ms):

	gcc 9.2.0	icc 19.0.4.243	
	glibc 2.17	gcc version 9.2.0 compat.	
single	0.5	0.4 [1]	
double	1	0.5	
quadruple	62 [2]	10 [3]	

[1] results differ with optimization level 0 ($x_2 = -0.654694$, $y_2 = 0.631660$), level 1 ($x_2 = -1.343469$, $y_2 = 0.625392$), and levels 2 or 3 ($x_2 = -1.182620$, $y_2 = 0.601759$) [2] time extrapolated on another machine

[3] compiled with -Qoption, cpp, -extended_float_types

What About Accuracy?

Tested with mpcheck (mpcheck.gforge.inria.fr) based on GNU MPFR. 10⁶ random tests. Rounding to nearest. Error in ulps.

function	glibc 2.31	icc 19.0.4.243
exp	0.501	0.501
log	0.871	0.501
log2	2.14	0.501
log10	1.43	0.501
sin	1.27	0.501
atan	1.09	0.501
acos	1.13	0.501
sinh	1.83	0.501
tanh	2.30	0.501
acosh	3.24	0.501
tgamma	4.70	4090 [1]

[1] bug reported, for $x = 0 \times 3.08 \text{e} 1 \text{f} 38 \text{d} \text{d} 769117414 \text{b} \text{f} 11 \text{b} 45 \text{d} \text{c} \text{p} + 8$

How Slow is Quadruple Precision?

If we replace all calls to sinf128 by the following (same for cosf128):

```
static _Float128 my_sinf128 (_Float128 x)
{
   return (_Float128) 0.5;
}
```

the total time is divided by 18.1 with glibc, by 7.1 with the Intel Math Library.

Conclusion: the main bottleneck are the mathematical functions.

On our double pendulum example, quadruple precision is 20 times slower than double precision with the Intel Math Library, and 62 times with the GNU library.

Can we do better?

Challenge: implement a fast exp function in quadruple precision.

Target processor: x86_64.

The GNU libm takes on average about 3200 cycles.

The Intel Math Library takes on average 280-430 cycles.

Goal: save a factor of 10 over the GNU libm.

Everything is allowed.

Accuracy constraint: should be about as accurate as the glibc function.

Time constraint: at most one week of design/coding/testing (March 23-27, 2020).

Principle 1: avoid all operations on _Float128, even addition and multiplication.

Instead, extract the _Float128 input into a special binary128 structure, do all computations on binary128, and unpack at the end.

The binary128 structure

Encoding:

sign s	exponent <i>e</i>	m_0	m_1
int	int	uint64_t	uint64_t
$m{s} \in \{-1,1\}$	$-16493 \le e \le 16383$	$m_0 < 2^{64}$	$m_1 < 2^{64}$

Decoding:

$$x = s \cdot 2^{e} \cdot \left(\frac{m_1}{2^{64}} + \frac{m_0}{2^{128}}\right)$$

Encoding similar to GNU MPFR, with no implicit bit.

No systematic normalization (m_1 can be smaller than 2^{63}).

Corollary: we get 128 - 113 = 15 extra bits of accuracy.

How Slow is Quadruple Precision?

Algorithm for binary128 exponential

- extract x into a binary128 structure, say y
- 2. check for special values, overflow, underflow
- 3. write $y = i \log 2 + i \log 2 \cdot 2^{-8} + k \log 2 \cdot 2^{-16} + r$ with $-128 \le i, k \le 128$ and $|r| \le \log 2 \cdot 2^{-17}$ 4. $v \leftarrow v - i \log 2$ $u_i \approx j \log 2 \cdot 2^{-8}, v_k \approx k \log 2 \cdot 2^{-16}$ 5. $v \leftarrow v - u_i - v_k$ $f_i \approx \exp(u_i), g_k \approx \exp(v_k)$ 6. $e_{ik} \leftarrow f_i \cdot g_k$ 7. now $|y| \le \log 2 \cdot 2^{-17}$ 8. $z \leftarrow v(p_4 + v(p_5v + p_6))$ [64-bit arithmetic only] 9. $z \leftarrow p_1 + v(p_2 + v(p_3 + z))$ 10. $\mathbf{y} \leftarrow \mathbf{e}_{ik} + \mathbf{y} \cdot \mathbf{z} \cdot \mathbf{e}_{ik}$ 11. return unpack(y, i) [multiplies by 2^{i}]

The coefficients $p_1, p_2, ..., p_6$ were generated by the Sollya tool.

$$p(x) = 1 + p_1 x + p_2 x^2 + \dots + p_6 x^6$$

They minimize the relative error of $p(x) - \exp x$ for $|x| \le \log 2 \cdot 2^{-17}$, with the following constraints: p_1, p_2, p_3 fit on 128 bits p_4, p_5, p_6 fit on 64 bits

- $p_5 = 0 \times 2.22222224 \text{dce8p-8}$

$$p_6 = 0 \times 5.b05b43776501cp-12$$

Relative error $< 2^{-121.33}$ (not counting rounding errors).

How Slow is Quadruple Precision?

[a, b, c stand for binary128 structures, m stands for some $m_1 \cdot 2^{-64} + m_0 \cdot 2^{-128}$ with $m_1, m_0 < 2^{64}$ extract_binary128: extract a _Float128 into binary128 unpack: unpack a binary128 into a _Float128 normalize: shift m_1, m_0 and adjust e so that $2^{63} < m_1 < 2^{64}$ align_binary128: shift *m* so that e = 0 (assumes $e \le 0$ initially) sub_inplace: $a \leftarrow a - c$, assuming $e_a = e_c$ add inplace: $a \leftarrow a + c$ mul: $a \leftarrow \operatorname{high}(b \cdot c)$ addu: $a \leftarrow b + m \cdot 2^{e_b}$, assuming no carrv shift right, shift left: shift m_a by k bits and update e_a

reduce: $a \leftarrow a - i \log 2$, *i* integer, log 2 precomputed on 192 bits

Test done by Alexei Sibidanov, on 10^5 random inputs in [-10, 10].

Correctly rounded results:

Oracle	Intel Math	libquadmath	ASquadmath	Paul's exp
Studio 12.6	19.0.5.281	9.2.1		
99615	99997	99999	99999	99951

All other results are wrong by one ulp.

Test done by Alexei Sibidanov, on an AMD Ryzen 5 2400G.

Average number of cycles (measured with perf stat):

MPFR	Oracle Studio	Intel Math	libquad-	ASquad-	Paul's
4.0.2	12.6	19.0.5.281	math 9.2.1	math	exp
6213	7634	427	3142	181	234

Goal of saving a factor of 10 over the GNU libm is reached!

Performance of exp function



[credit Alexei Sibidanov]

Conclusion

Quadruple precision is indeed slow, but we can do much better!

We saved a factor of 10 with little effort, probably we can save another factor of 2 with more effort.

Use of integer operations is the key for efficiency.

The generic binary128 routines can be reused for other functions.

Implement addition, subtraction, multiplication, division directly with correct rounding for the binary128 type, to avoid converting to/from _Float128.

Design an exponential function with correct rounding. The slow path would use similar integer-only arithmetic, with four 64-bit words (256 bits of accuracy), assuming the hard-to-round cases are known.

The libm detector

https://homepages.loria.fr/PZimmermann/libm-detector/

```
$ gcc libm-detector.c -lm
$ ./a.out
Mathematical Library Detector, version 1.0
Probably libm shipped with GNU libc, version >= 2.29
```

```
$ icc libm-detector.c
$ ./a.out
Mathematical Library Detector, version 1.0
Probably Intel Math Library
```

Performance is nice.

What about reproducibility (cf. David Bailey's talk)?

Currently developers/users mostly care about efficiency.

Shouldn't we instead seek for bit-to-bit reproducibility, and then only for performance?

IEEE 754 only recommends correct rounding for math functions. Should it require correct rouding?

Maybe we get some answer in Jason Riedy's talk Potential Directions for Moving IEEE-754 Forward at 3:45pm.

Sollya: an environment for the development of numerical codes, Sylvain Chevillard, Mioara Maria Joldes and Christoph Lauter, Third International Congress on Mathematical Software (ICMS), LNCS 6327, 2010.

A new quadruple precision math library, Alexei Sibidanov, 50 pages, personal communication, February 2020.

Source code for expf128 available here:

https://homepages.loria.fr/PZimmermann/glibc-contrib/